

# A note on light geometric graphs

Eyal Ackerman\*

Jacob Fox†

Rom Pinchasi‡

February 16, 2012

## Abstract

Let  $G$  be a geometric graph on  $n$  vertices in general position in the plane. We say that  $G$  is *k-light* if no edge  $e$  of  $G$  has the property that each of the two open half-planes bounded by the line through  $e$  contains more than  $k$  edges of  $G$ . We extend the previous result in [1] and with a shorter argument show that every  $k$ -light geometric graph on  $n$  vertices has at most  $O(n\sqrt{k})$  edges. This bound is best possible.

**Keywords:** Geometric graphs,  $k$ -near bipartite.

## 1 Introduction

Let  $G$  be an  $n$ -vertex *geometric graph*. That is,  $G$  is a graph drawn in the plane such that its vertices are distinct points and its edges are straight-line segments connecting corresponding vertices. It is usually assumed, as we will assume in this paper, that the set of vertices of  $G$  is in general position in the sense that no three of them lie on a line.

A typical question in geometric graph theory asks for the maximum number of edges that a geometric graph on  $n$  vertices can have assuming a forbidden configuration in that graph. This is a popular area of study extending classical extremal graph theory, utilizing diverse tools from both geometry and combinatorics. For example, an old result of Hopf and Pannwitz [3] and independently Sutherland [7] states that any geometric graph on  $n$  vertices with no pair of disjoint edges has at most  $n$  edges. This is a special case of Conway's thrackle conjecture.

Let  $e$  be an edge of  $G$ . We say that  $G$  has a *k-light side* with respect to  $e$ , if one of the two open half-planes bounded by the line through  $e$  contains at most  $k$  edges of  $G$ . If  $G$  has a *k-light side* with respect to *every* edge  $e$ , then we say that  $G$  is *k-light*. In other words,  $G$  is  $k$ -light if no edge of  $G$  has the property that each of the two open half-planes bounded by the line through  $e$  contains more than  $k$  edges of  $G$ .

The notion of a  $k$ -light graph is a weakening of the notion of a  $k$ -near bipartite graph defined in [1]. A graph  $G$  is  $k$ -near bipartite if every line in the plane bounds an open half plane containing at most  $k$  edges of  $G$ . Therefore, every  $k$ -near bipartite graph is also a  $k$ -light graph. It is shown in [1] that  $k$ -near bipartite graphs on  $n$  vertices contain  $O(\sqrt{kn})$  edges. In this paper we prove the same result for  $k$ -light graphs, thus strengthening the result in [1]. Moreover, our proof is much shorter but on the other hand relies on other results about geometric graphs.

---

\*Department of Mathematics, Physics, and Computer Science, University of Haifa at Oranim, Tivon 36006, Israel. [ackerman@sci.haifa.ac.il](mailto:ackerman@sci.haifa.ac.il).

†

‡Mathematics Department, Technion—Israel Institute of Technology, Haifa 32000, Israel. [room@math.technion.ac.il](mailto:room@math.technion.ac.il). Supported by BSF grant (grant No. 2008290).

## 2 The maximum number of edges in $k$ -light geometric graphs

We are interested in the maximum number of edges of an  $n$ -vertex  $k$ -light geometric graph. A simple construction from [1] shows an  $n\sqrt{k}$  lower bound for  $k \leq (\frac{n}{2} - 1)^2$ , even for  $k$ -near bipartite graphs. In this construction every line contains at most  $k$  edges of  $G$  in one of the two open half-planes bounded by it. Another construction of a  $k$ -light graph with  $n\sqrt{k}$  edges is obtained by taking the vertices of a regular  $n$ -gon and connecting by edges vertices whose cyclic distance is at most  $\sqrt{k}$ . In this construction, however, it is no longer true that every line bounds an open half-plane containing at most  $k$  edges of  $G$ .

Our main result shows that these constructions are essentially best possible.

**Theorem 1.** *Let  $n$  and  $k$  be positive integers. Every  $n$ -vertex  $k$ -light geometric graph has at most  $O(n\sqrt{k})$  edges.*

**Proof.** Let  $G$  be an  $n$ -vertex  $k$ -light geometric graph with  $m$  edges. We orient every edge  $e$  of  $G$  in such a way that the open half-plane bounded to the left of  $e$  contains at most  $k$  edges of  $G$ . Because  $G$  is  $k$ -light such an orientation exists.

We will need the following two lemmas.

**Lemma 2.1.** *Let  $G$  be an oriented geometric graph on  $n$  vertices. There exists an absolute constant  $c_3$  such that if  $G$  has more than  $c_3n$  edges, then it contains an edge  $e$  such that the open half-plane bounded to the left of  $e$  contains an edge of  $G$ .*

**Proof.** It is enough to show that in any (unoriented) geometric graph  $G$  with  $n$  vertices and sufficiently many (that is, at least  $c_3n$ ) edges there is an edge  $e$  such that each of the two open half-planes bounded by the line through  $e$  contains an edge of  $G$ . This is in fact the case  $k = 1$  in Theorem 1 that we wish to prove. The reader is encouraged to find a simple proof of this fact. Here we will rely on a rather elaborate argument of Valtr [8] that proves a much stronger statement than what we need.

We refer the reader to [5, 4, 8]. Two edges of a geometric graph are called *avoiding* or sometimes *parallel* if no line passing through one edge meets the other edge. Equivalently, two edges are avoiding if they are opposite edges in a convex quadrilateral.

The notion of avoiding edges was first defined by Kupitz [5], who conjectured that any geometric graph on  $n$  vertices with more than  $2n - 2$  edges must contain a pair of avoiding edges. In [4] it is shown that if a graph  $G$  on  $n$  vertices does not contain a pair of avoiding edges, then the number of edges in  $G$  is at most  $2n - 1$ . In [8] Valtr improved this bound by one, completing the proof of Kupitz' conjecture. He further generalized this result, showing that for any fixed  $k$ , every geometric graph with more than  $c_k n$  edges contains  $k$  pairwise avoiding edges. Here  $c_k$  is an absolute constant that depends only on  $k$ .

In fact, Valtr's result is a bit stronger. Looking into the proof in [8] reveals that he actually shows that a geometric graph with more than  $c_k n$  edges contains  $k$  edges  $e_1, \dots, e_k$  that are pairwise avoiding, but what is more important to our needs is that the line through  $e_i$  separates  $e_1, \dots, e_{i-1}$  from  $e_{i+1}, \dots, e_k$ . More specifically, Valtr defines three partial orders on a set of edges in  $G$  and any chain with respect to any of the partial orders is a collection of such edges. It is then shown that if the number of edges in  $G$  is large enough, then there exists a chain of length  $k$  in one of the partial orders.

Thus, for the case  $k = 3$  it follows that if  $G$  contains more than  $c_3 n$  edges, then there are three pairwise avoiding edges  $e, f, g$  such that the line through  $f$  separates  $e$  and  $g$ . This immediately implies Lemma 2.1, as in any orientation of  $f$  the half-plane bounded to the left of  $f$  will contain an edge of  $G$ . ■

**Lemma 2.2.** *Let  $G$  be an oriented geometric graph on  $n$  vertices with  $m$  edges. There exists a positive absolute constant  $d$  with the following property. If the number of edges in  $G$  is greater than  $2c_3n$  (where  $c_3$  is the constant from Lemma 2.1), then  $G$  contains at least  $dm^3/n^2$  pairs of edges  $(e, f)$  such that the open half-plane bounded to the left of  $e$  contains  $f$ .*

**Proof.** This is by now a quite standard consequence of the result in Lemma 2.1 and is carried out by a similar probabilistic technique used to derive a similar bound for the number of pairs of crossing edges in a geometric graph (see p. 55 in [6], also p. 45 in [2]).

Denote by  $x(G)$  the number of pairs of edges  $(e, f)$  in  $G$  such that the open half-plane bounded to the left of  $e$  contains  $f$ . Pick every vertex of  $G$  independently with probability  $p$ , and denote by  $G' = (V', E')$  the subgraph of  $G$  that is induced by the chosen vertices. Clearly,  $\mathbb{E}[|V'|] = pn$ ,  $\mathbb{E}[|E'|] = p^2m$ , and  $\mathbb{E}[x(G')] = p^4x(G)$ . On the other hand, it follows from Lemma 2.1 that  $x(G') \geq |E'| - c_3|V'|$ , and this holds also for the expected values:  $\mathbb{E}[x(G')] \geq \mathbb{E}[|E'|] - c_3\mathbb{E}[|V'|]$ . Plugging in the expected values and setting  $p = 2c_3n/m < 1$  we get that  $x(G) \geq \frac{1}{8c_3^2} \frac{m^3}{n^2}$ . ■

Let  $c_3$  and  $d$  be the constants from Lemmas 2.1 and 2.2. Clearly we may assume that  $G$  contains at least  $2c_3n$  edges or else we are done. By Lemma 2.2,  $G$  contains at least  $dm^3/n^2$  pairs  $(e, f)$  of edges such that the open half-plane bounded to the left of  $e$  contains  $f$ . However, by the choice of orientation of the edges in  $G$ , an edge  $e$  can belong to at most  $k$  such pairs  $(e, f)$ . We conclude that  $dm^3/n^2 \leq km$ . This now easily implies that  $m \leq \frac{1}{\sqrt{d}}n\sqrt{k}$  as desired. ■

## References

- [1] E. Ackerman and R. Pinchasi, On the light side of geometric graphs, *Discrete Math.*, in press.
- [2] S. Felsner, Geometric graphs and arrangements. Some chapters from combinatorial geometry. Advanced Lectures in Mathematics. Friedr. Vieweg & Sohn, Wiesbaden, 2004.
- [3] H. Hopf and E. Pannwitz, Aufgabe Nr. 167, *Jahresbericht. Deutsch. Math.-Verein.* **43** (1934), 114.
- [4] M. Katchalski and H. Last, On geometric graphs with no two edges in convex position, *Discrete Comput. Geom.* **19** (1998), no. 3, Special Issue, 399–404.
- [5] Y. Kupitz, On pairs of disjoint segments in convex position in the plane, *Ann. Discrete Math.* **20** (1984), 203–208.
- [6] J. Matoušek, Lectures on discrete geometry. Graduate Texts in Mathematics, 212. Springer-Verlag, New York, 2002.
- [7] J. W. Sutherland, Lösung der Aufgabe 167, *Jahresbericht Deutsch. Math.-Verein.* **45** (1935), 33–35.
- [8] P. Valtr, On geometric graphs with no  $k$  pairwise parallel edges, *Discrete Comput. Geom.* **19** (1998), no. 3, 461–469.